

*On Multiplicity of Eigenvalues
in Robust Compliance Optimization
under Uncertain Loads*

Yoshihiro Kanno

(Tokyo Institute of Technology)

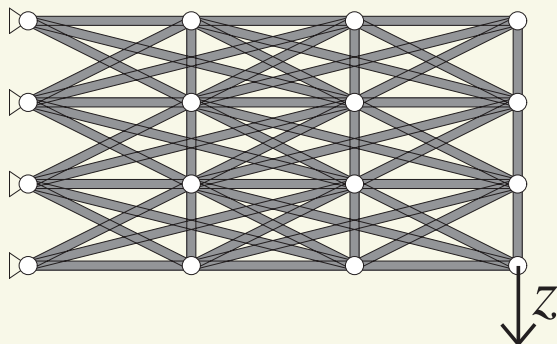
November 23, 2015 (ICMR 2015)

robust optimization

- nominal (i.e., conventional) compliance optimization

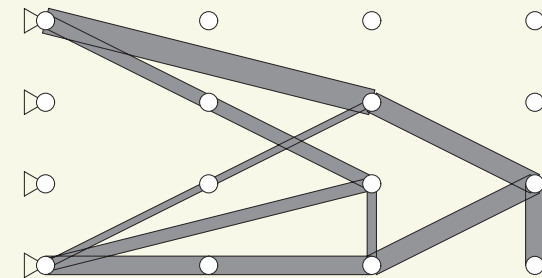
$$\begin{aligned} \min \quad & \text{compl}(\mathbf{x}; \mathbf{z}) \\ \text{s. t.} \quad & \mathbf{x} \geq \mathbf{0}, \quad \text{vol}(\mathbf{x}) \leq \bar{V} \end{aligned}$$

- \mathbf{x} : cross-sectional areas \mathbf{z} : data (e.g., external load)



ground struct.

→
optimize



nominal opt. sol.

robust optimization

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- robust compliance optimization

$$\begin{aligned} \min \quad & \max\{\text{compl}(\mathbf{x}; \mathbf{z}) \mid \mathbf{z} \in \mathcal{U}\} \\ \text{s. t.} \quad & \mathbf{x} \geq \mathbf{0}, \quad \text{vol}(\mathbf{x}) \leq \bar{V} \end{aligned}$$

- \mathcal{U} : uncertainty set (e.g., set of uncertain loads)
- obj. fcn.: worst value of compliance

robust optimization

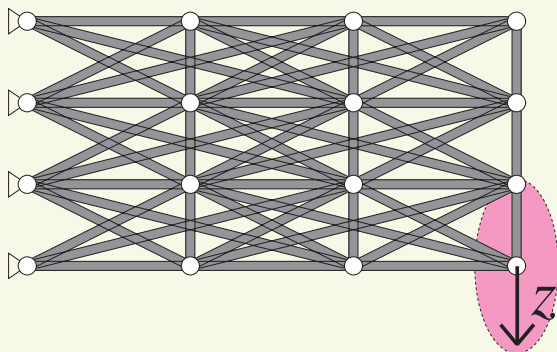
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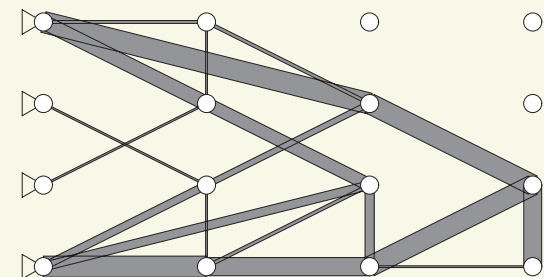
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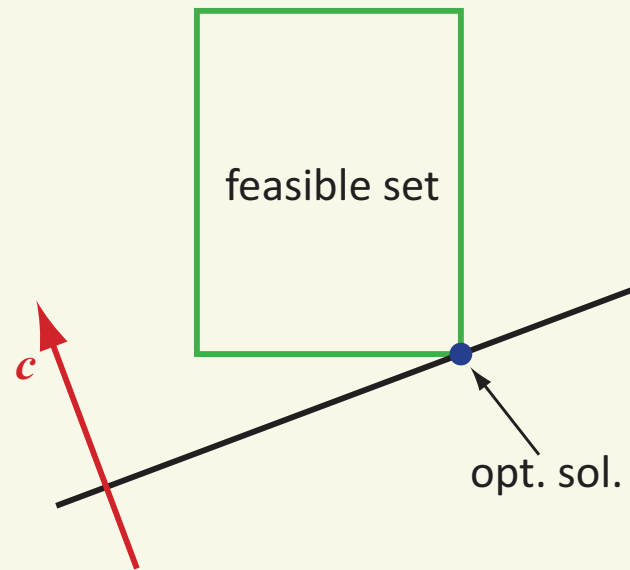
uncertainty: $\mathbf{z} \in \mathcal{U}$

→
optimize



robust opt. sol.

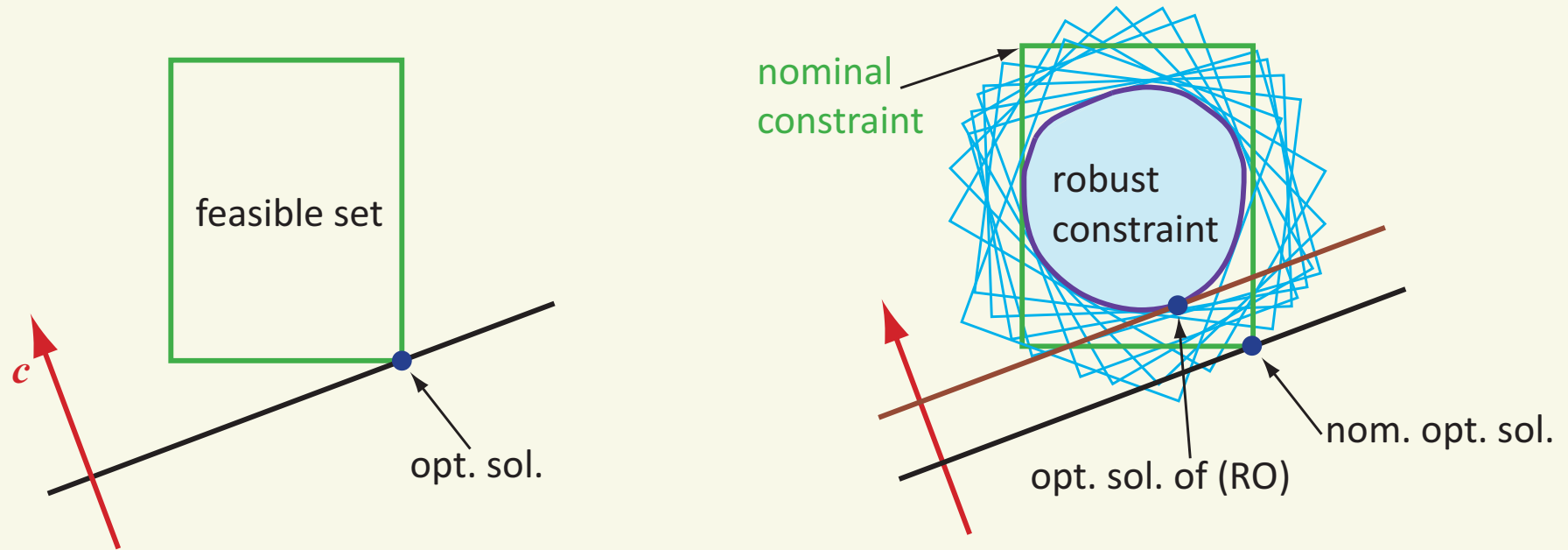
robust optimization



- nominal opt.:

$$\begin{aligned} \min \quad & \mathbf{c}^\top \mathbf{x} \\ \text{s. t.} \quad & l_i \leq x_i \leq u_i \end{aligned}$$

robust optimization



- nominal opt.:

$$\begin{aligned} \min \quad & \mathbf{c}^\top \mathbf{x} \\ \text{s. t.} \quad & l_i \leq x_i \leq u_i \end{aligned}$$

- robust opt.:

$$\begin{aligned} \min \quad & \mathbf{c}^\top \mathbf{x} \\ \text{s. t.} \quad & l_i(\mathbf{z}) \leq x_i \leq u_i(\mathbf{z}) \quad (\forall \mathbf{z} \in \mathcal{U}) \end{aligned} \quad \text{(RO)}$$

SDPs for robust optimization

- SDP = semidefinite programming — convex
 - compliance, load unc. [Ben-Tal & Nemirovski '97]
 - compliance, geometric unc. (safe approx.) [Hashimoto & K. '15]
- nonlinear SDP — nonconvex
 - stress cstr., load unc. [K. & Takewaki '06]
 - stiffness unc. (safe approx.) [Guo, Bai, Zhang, & Gao '09]
 - stiffness & length unc. (safe approx.) [Guo, Du, & Gao '11]

fundamental from linear algebra

- def.

$B - C^T A^{-1} C$: the Schur complement of A in $\left[\begin{array}{c|c} A & C \\ \hline C^T & B \end{array} \right]$

- A, B : symm. matrices

fundamental from linear algebra

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$$B - C^T A^{-1} C : \text{ the Schur complement of } A \text{ in } \left[\begin{array}{c|c} A & C \\ \hline C^T & B \end{array} \right]$$

- A, B : symm. matrices
- Lemma on the Schur complement
 - Assume A is p.d.
 - Then,

$$\left[\begin{array}{c|c} A & C \\ \hline C^T & B \end{array} \right] \text{ is p.s.d.} \iff B - C^T A^{-1} C \text{ is p.s.d.}$$

-
- p.d. = positive definite \iff all eigenvalues > 0 .
 - p.s.d. = positive semidefinite \iff all eigenvalues ≥ 0 .

problem setting

- f : external load
- uncertainty set

$$\mathcal{U} = \{Qe \mid \|e\| \leq 1\}$$

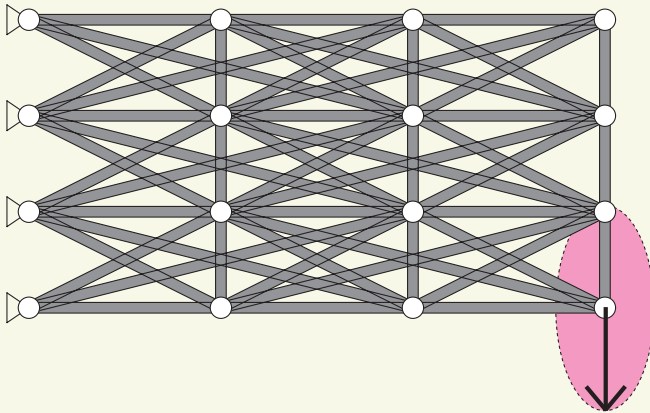
- Q : constant matrix
 - info. of “nom. load & unc. load”

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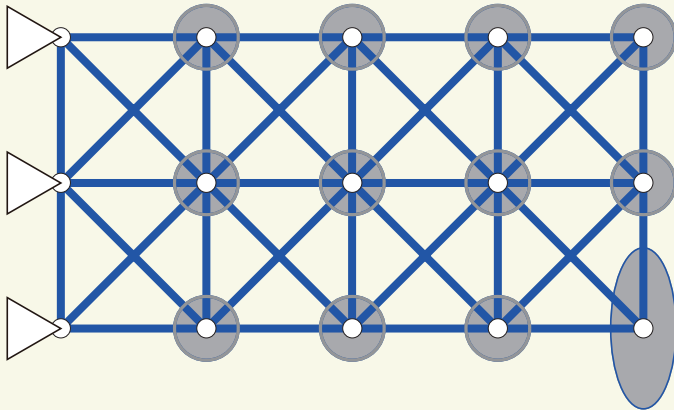
$$Q = \left[\begin{array}{c|c} \tilde{f} & 0 \\ 0 & \epsilon \\ \hline 0 & 0 \\ \vdots & \vdots \\ 0 & 0 \end{array} \right], \quad e = \begin{bmatrix} e_1 \\ e_2 \end{bmatrix}$$

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$$Q = \begin{bmatrix} \tilde{f} & 0 & 0 & \cdots & 0 \\ 0 & \epsilon & 0 & \cdots & 0 \\ 0 & 0 & \epsilon & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & \epsilon \end{bmatrix}, \quad e = \begin{bmatrix} e_1 \\ e_2 \\ e_3 \\ \vdots \\ e_d \end{bmatrix}$$

problem setting

- f : external load
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- Q : constant matrix
 - info. of “nom. load & unc. load”
- RO:

$$\begin{aligned} \min_x \quad & \max_f \{ \text{compl}(\mathbf{x}; f) \mid f \in \mathcal{U} \} \\ \text{s. t.} \quad & \mathbf{x} \geq \mathbf{0}, \quad \text{vol}(\mathbf{x}) \leq \bar{V} \end{aligned}$$

- \mathbf{x} : c.-s. areas — design variables
- Tractable reformulations are known.

tractable forms of (RO)

- three reformulations in literature

- SDP formulation

[Ben-Tal & Nemirovski '97]

- min. of maximum eigenvalue (standard eig. prob.)

[Takezawa, Nii, Kitamura, & Kogiso '11]

- min. of maximum eigenvalue (generalized eig. prob.)

[Cherkaev & Cherkaev '03, '08]

- equivalence \leftarrow the Schur complement lemma

SDP to standard eig. prob.

- SDP formulation

[Ben-Tal & Nemirovski '97]

$w \geq$ (worst compl.)

$$\Leftrightarrow \left[\begin{array}{c|c} K(\mathbf{x}) & Q \\ \hline Q^\top & wI \end{array} \right] \text{ is p.s.d.} \quad (\spadesuit)$$

- $K(\mathbf{x})$: stiffness matrix
- w : upr. bd. for worst compl.

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- Apply the Schur complement lemma (assuming $K(\mathbf{x})$ is invertible):

$$(\spadesuit) \Leftrightarrow wI - Q^\top K(\mathbf{x})^{-1} Q \text{ is p.s.d.}$$

$$\left[\begin{array}{c|c} A & C \\ \hline C^\top & B \end{array} \right] \text{ is p.s.d.} \quad \Leftrightarrow \quad B - C^\top A^{-1} C \text{ is p.s.d.}$$

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$$(\spadesuit) \Leftrightarrow wI - Q^\top K(\mathbf{x})^{-1} Q \text{ is p.s.d.}$$

- $\Leftrightarrow w \geq \lambda_{\max}(Q^\top K(\mathbf{x})^{-1} Q)$

- Therefore,

$$\text{(worst compl.)} = \lambda_{\max}(Q^\top K(\mathbf{x})^{-1} Q)$$

[Takezawa, Nii, Kitamura, & Kogiso '11]

SDP to generalized eig. prob.

- SDP formulation

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- Apply the Schur complement lemma (assuming $w > 0$):

$$(\spadesuit) \Leftrightarrow K(\mathbf{x}) - Q^\top (wI)^{-1} Q \text{ is p.s.d.}$$

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- Apply the Schur complement lemma (assuming $w > 0$):

$$(\spadesuit) \Leftrightarrow K(\mathbf{x}) - Q^\top (wI)^{-1} Q \text{ is p.s.d.}$$

- $\Leftrightarrow wK(\mathbf{x}) - QQ^\top$ is p.s.d. i.e., $\mathbf{z}^\top (wK(\mathbf{x}) - QQ^\top) \mathbf{z} \geq 0 \ (\forall \mathbf{z})$

- $\Leftrightarrow w \geq \frac{\mathbf{z}^\top QQ^\top \mathbf{z}}{\mathbf{z}^\top K(\mathbf{x}) \mathbf{z}} \ (\forall \mathbf{z})$ (Rayleigh–Ritz ratio)

- Therefore,

$$(\text{worst compl.}) = \max. \text{ eigenvalue of } (QQ^\top) \mathbf{z} = \lambda K(\mathbf{x}) \mathbf{z}$$

[Cherkaev & Cherkaev '03, '08]

- SDP formulation

[Ben-Tal & Nemirovski '97]

$$w \geq (\text{worst compl.}) \quad \Leftrightarrow \quad \left[\begin{array}{c|c} K(\mathbf{x}) & Q \\ \hline Q^\top & wI \end{array} \right] \text{ is p.s.d.} \quad (\spadesuit)$$

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- two ways of application of the Schur complement lemma

- (1): \rightarrow standard eig. prob. [Takezawa, Nii, Kitamura, & Kogiso '11]

$$(\spadesuit) \Leftrightarrow wI - Q^\top K(\mathbf{x})^{-1} Q \text{ is p.s.d.}$$

- (2): \rightarrow generalized eig. prob. [Cherkaev & Cherkaev '03, '08]

$$(\spadesuit) \Leftrightarrow K(\mathbf{x}) - Q^\top (wI)^{-1} Q \text{ is p.s.d.}$$

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on multiple eigenvalue

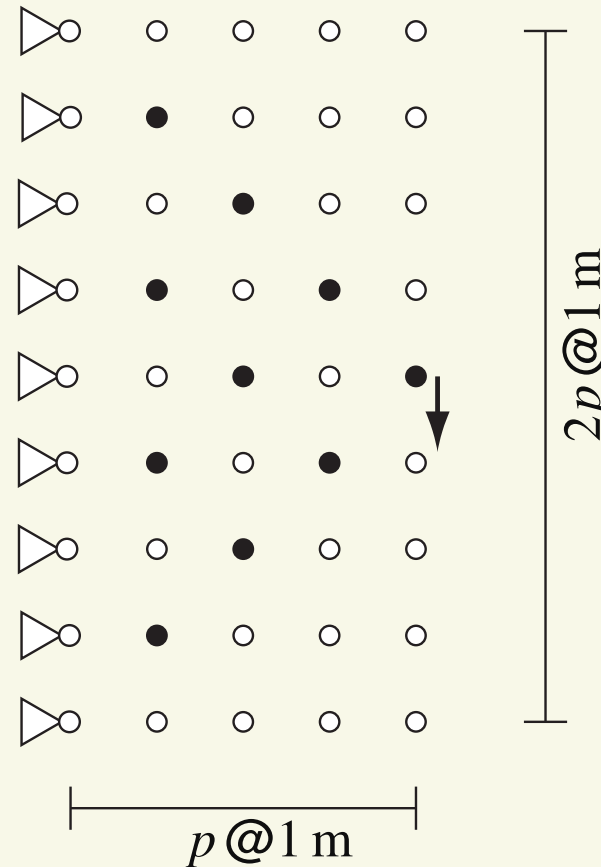
- the three formulations — related to eigenvalues
- eigenvalue optimization
 - In general, an opt. sol. often has multiple eigenvalues.
 - e.g., frequency, buckling load,...
 - ...How about (RO)?

on multiple eigenvalue

- the three formulations — related to eigenvalues
- eigenvalue optimization
 - In general, an opt. sol. often has multiple eigenvalues.
 - e.g., frequency, buckling load,...
 - ...How about (RO)? → Yes!
 - examples w/ simple eigenvalues
 - [Takezawa, Nii, Kitamura, & Kogiso '11]
 - [Brittain, Silva, & Tortorelli '12]
 - an example w/ fivefold eigenvalue
 - [Herskovits, Freire, Tanaka Fo, & Canelas '11]
 - ...but more.

a series of problem instances

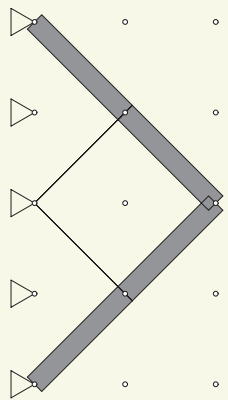
- truss (ground structure)



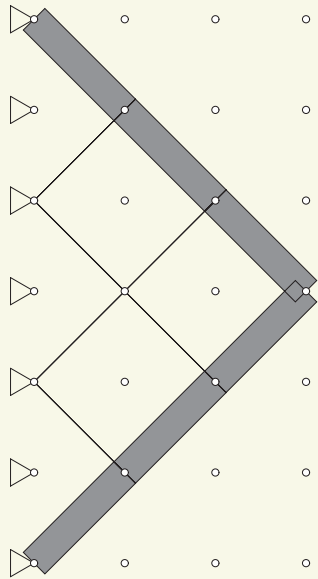
- Any two nodes are connected by a member.
 - Overlapping members have been removed.
- The problem size \rightarrow large, when $p \rightarrow$ large.

a series of problem instances

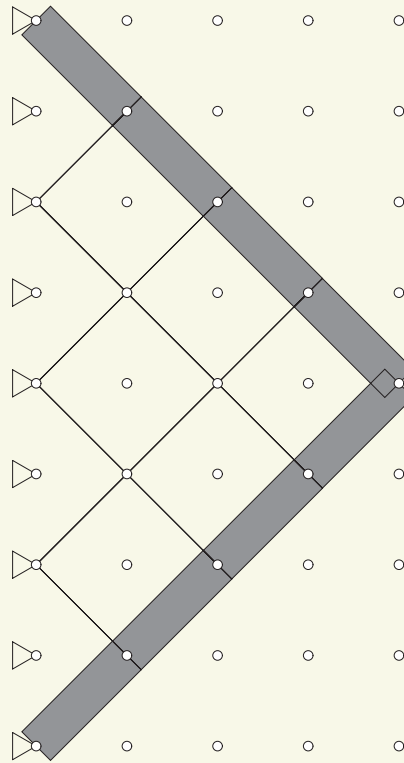
- robust optimal solutions:



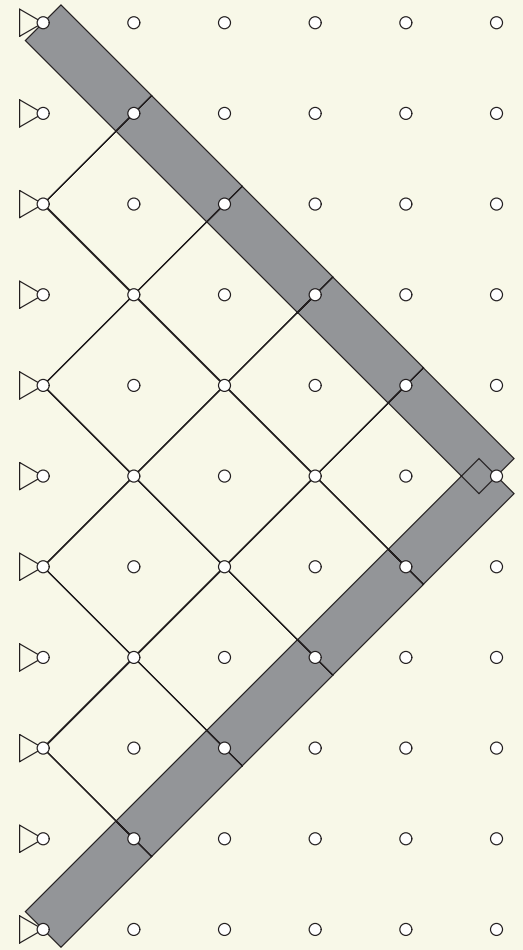
$p = 2$



$p = 3$



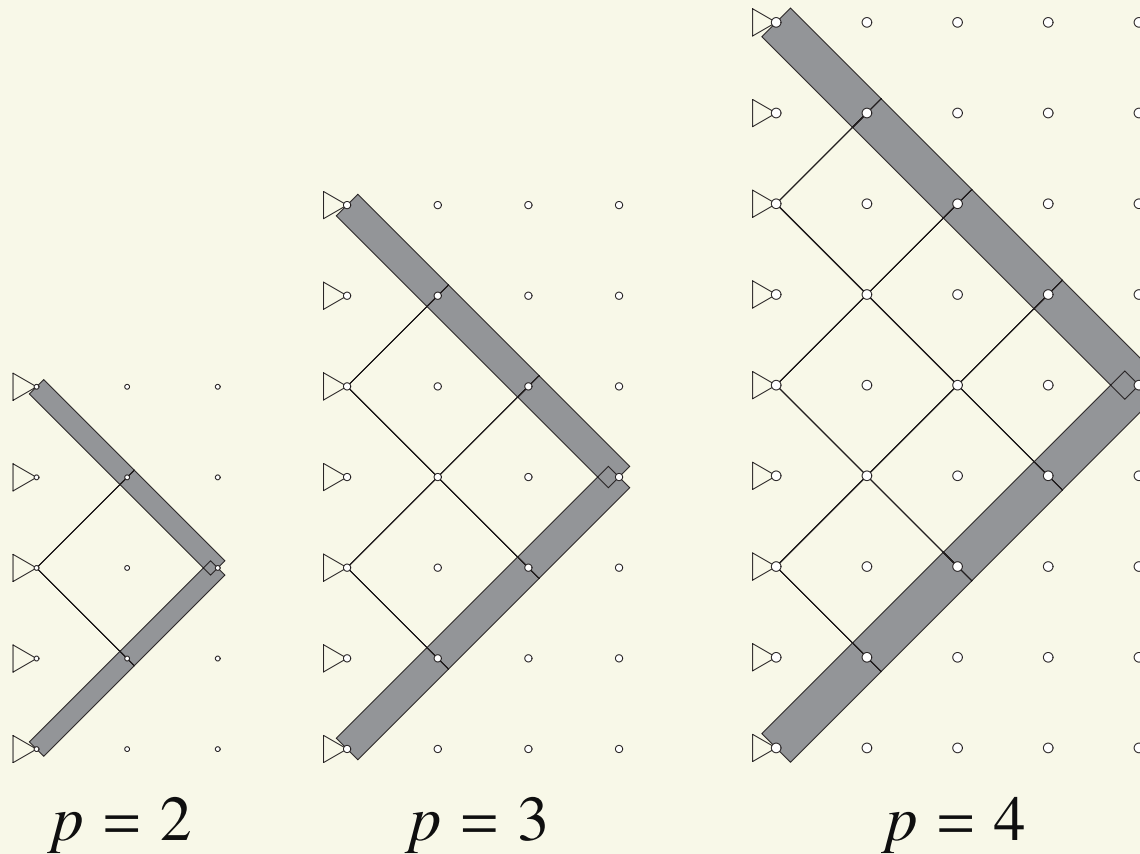
$p = 4$



$p = 5$

a series of problem instances

- robust optimal solutions:



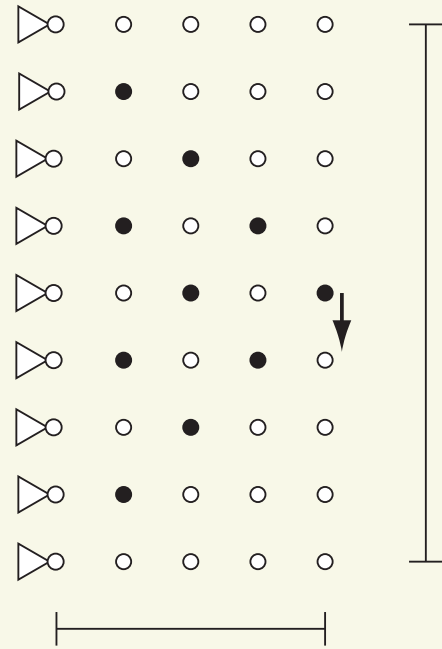
- multiplicity of eigenvalues = $2p - 1$

- As p increases, multiplicity increases.

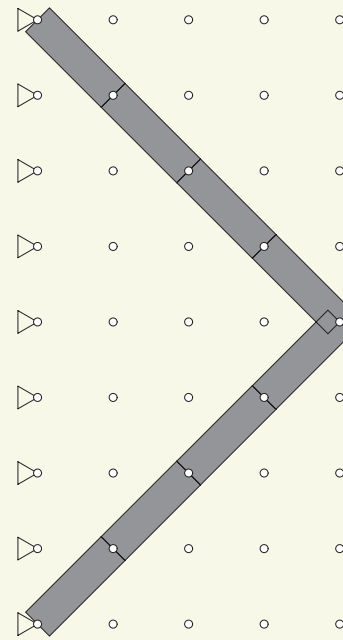
Arbitrary large multiplicity is possible.

source of multiplicity

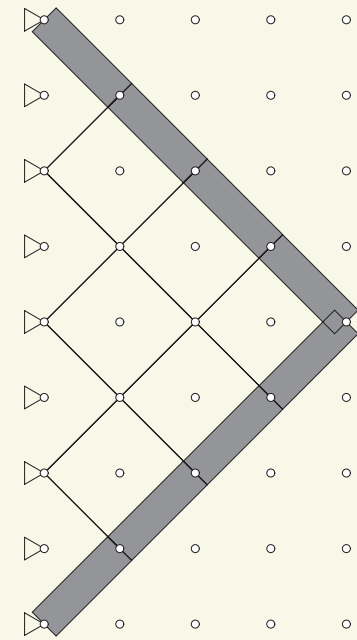
- $p = 4$



ground structure



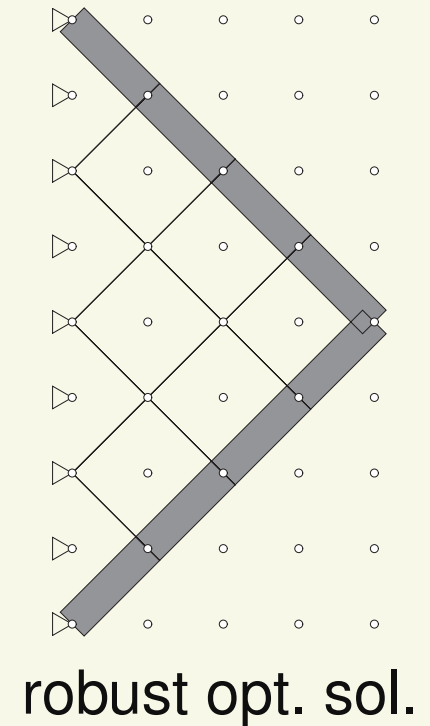
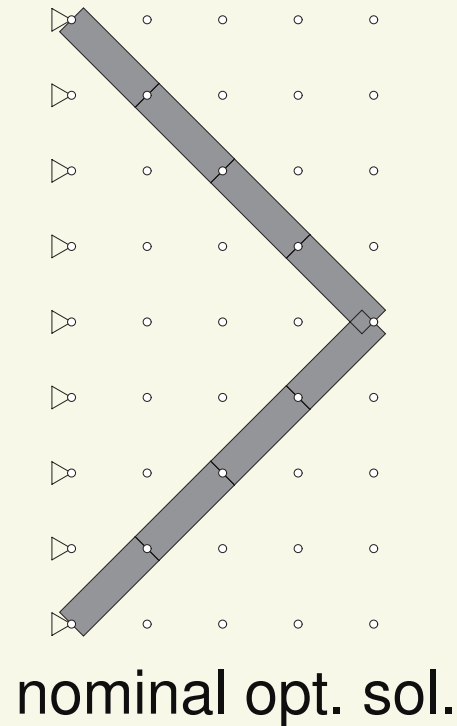
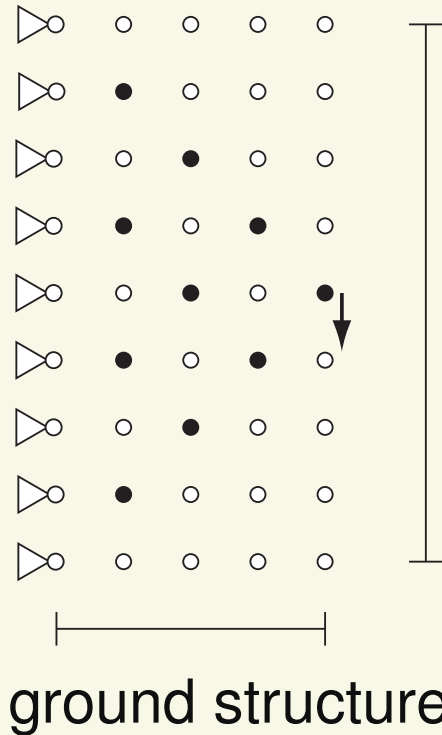
nominal opt. sol.



robust opt. sol.

source of multiplicity

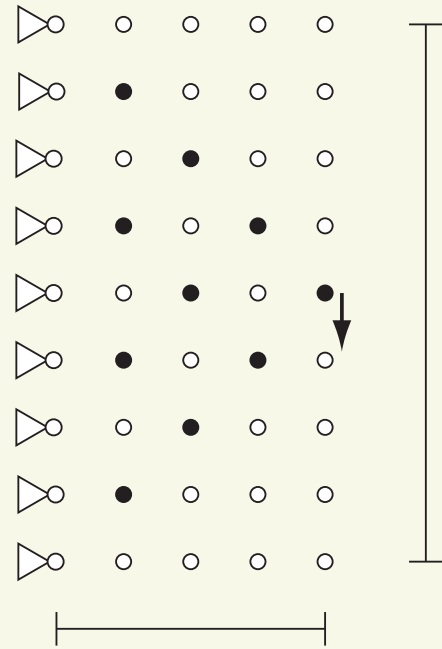
- $p = 4$



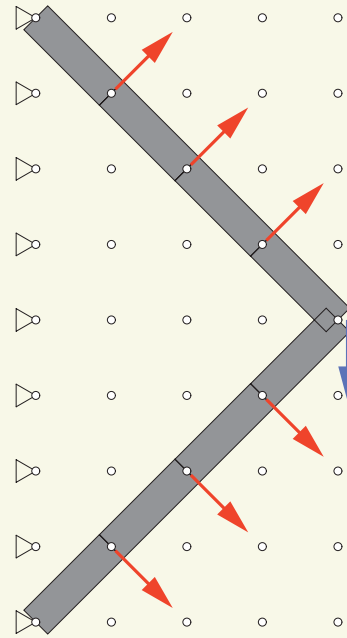
- multiple eigenvalues \Leftrightarrow worst-case loadings

source of multiplicity

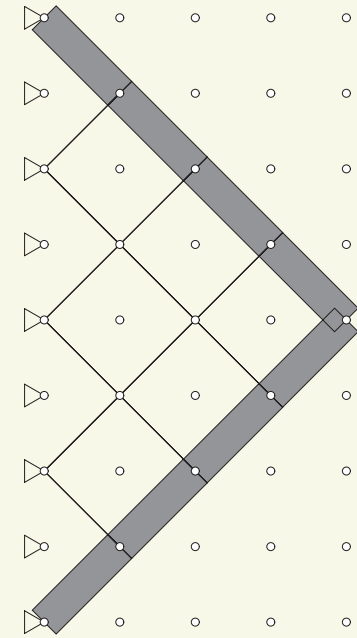
- $p = 4$



ground structure



nominal opt. sol.



robust opt. sol.

- multiple eigenvalues \Leftrightarrow worst-case loadings
 - 1 nominal load + $2(p - 1)$ uncertain loads @ unstable nodes
= $2p - 1$ multiplicity

conclusions

- robust compliance optimization
 - ellipsoidal uncertainty in external load
 - worst-case compliance → Min.
- three existing formulations
 - semidefinite programming [Ben-Tal & Nemirovski '97]
 - min. of max. eig. (standard eig. prob.) [Takezawa *et al.* '11]
 - min. of max. eig. (generalized eig. prob.) [Cherkaev & Cherkaev '03, '08]
 - → equivalence: explained via the Schur complement lemma
- multiplicity of eigenvalues at an opt. sol.
 - → a series of problems with arbitrary large multiplicity
 - multiple eigenvalues \Leftrightarrow multiple worst-case loadings
 - non-differentiability